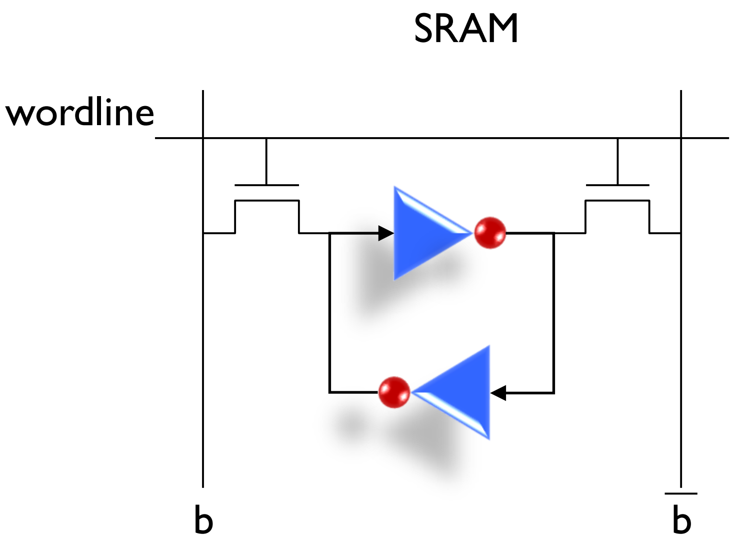
Main Memory and DRAM

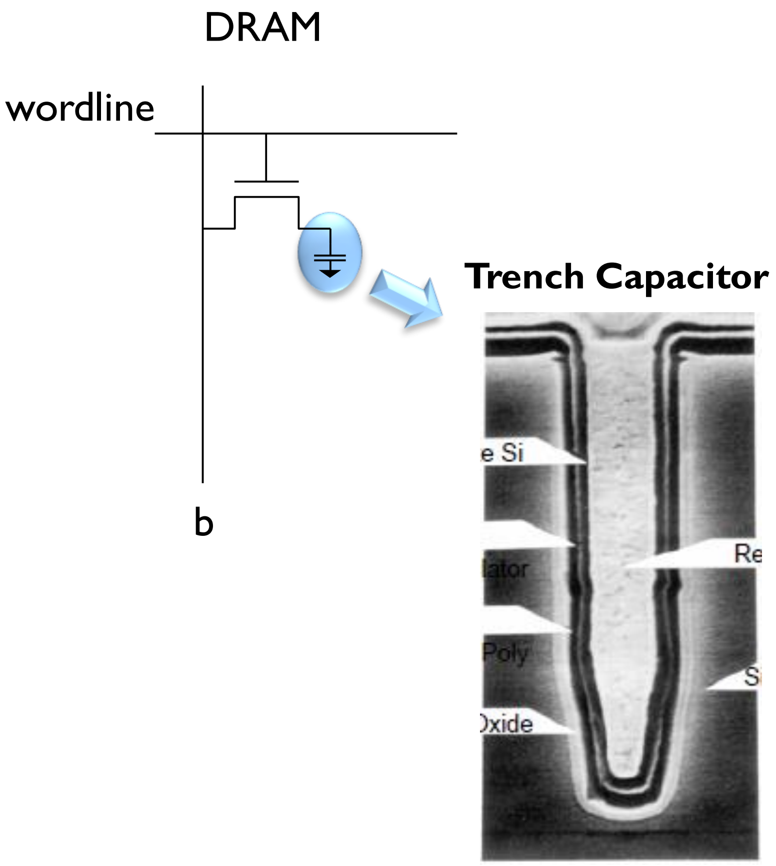
Instructor: Nima Honarmand

SRAM vs. DRAM

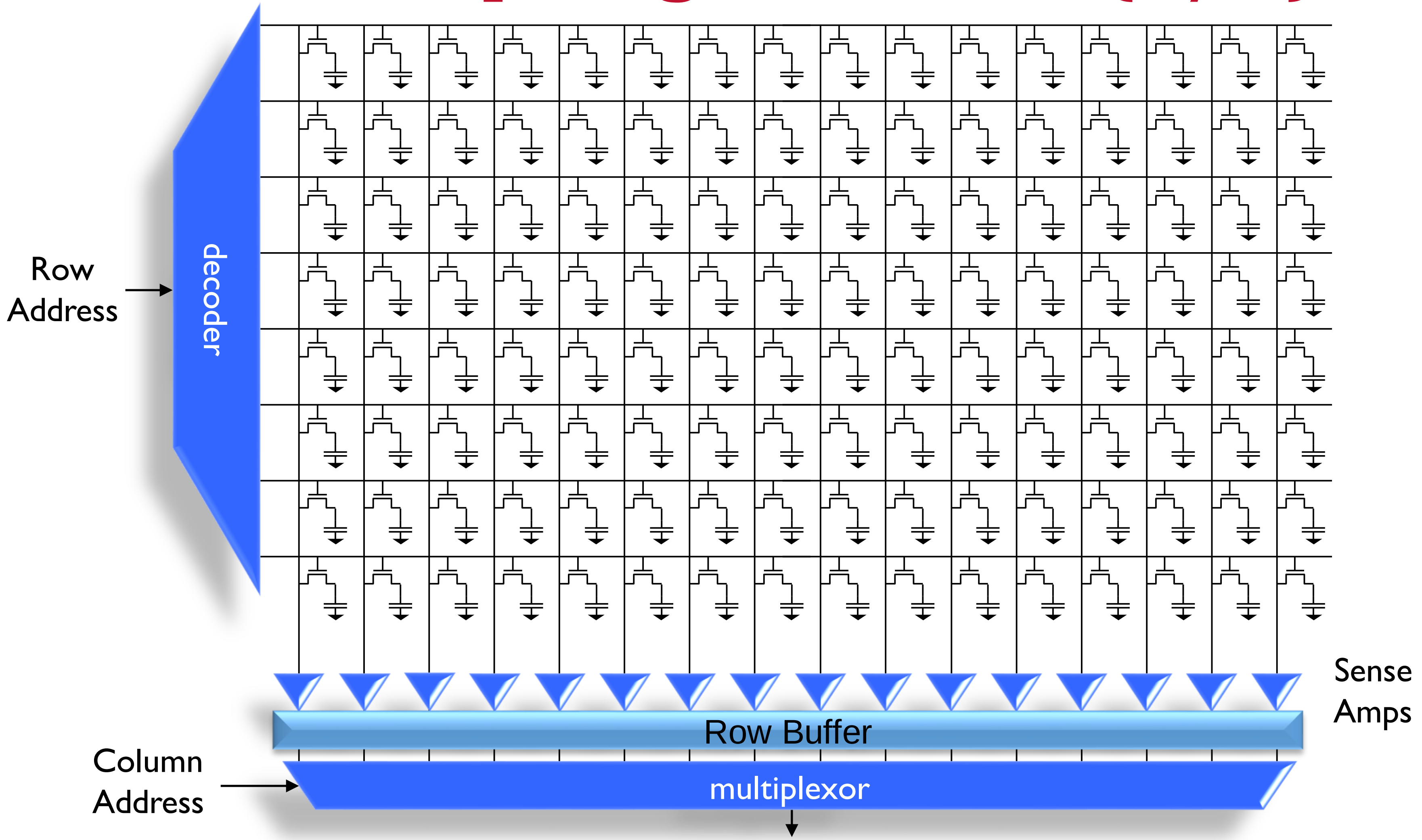
- SRAM = Static RAM
  - As long as power is present, data is retained
- DRAM = Dynamic RAM
  - If you don’t do anything, you lose the data
- SRAM: 6T per bit
  - built with normal high-speed CMOS technology
- DRAM: 1T per bit (+1 capacitor)
  - built with special DRAM process optimized for density

Hardware Structures





DRAM Chip Organization (1/2)

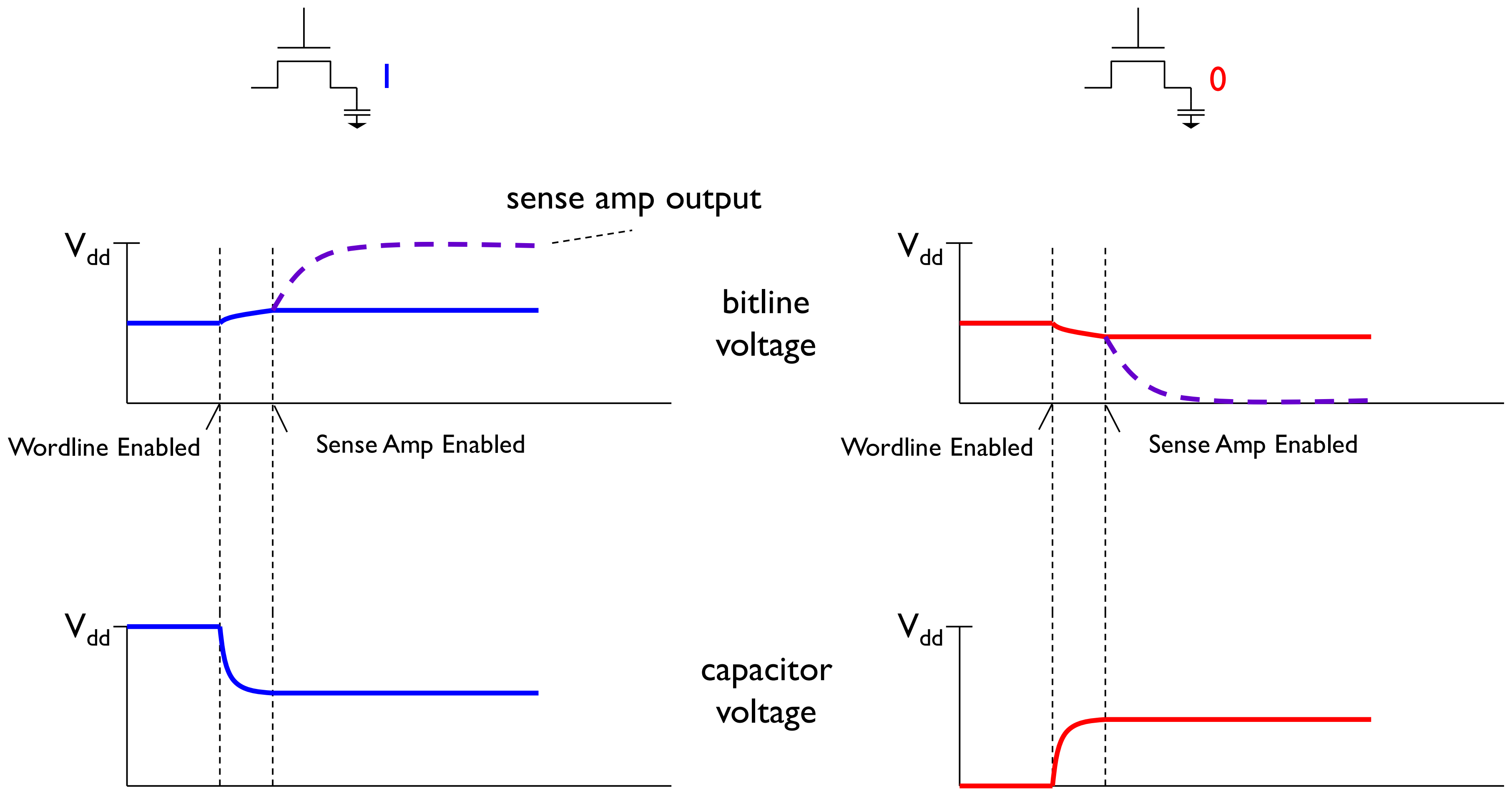


DRAM is much denser than SRAM

DRAM Chip Organization (2/2)

- Low-Level organization is very similar to SRAM
- Reads are *destructive*: contents are erased by reading
- *Row buffer* holds read data
  - Data in row buffer is called a *DRAM row*
    - Often called "page" – do not confuse with virtual memory page
  - Read gets entire row into the buffer
  - Block reads always performed out of the row buffer
    - Reading a whole row, but accessing one block
    - Similar to reading a cache line, but accessing one word

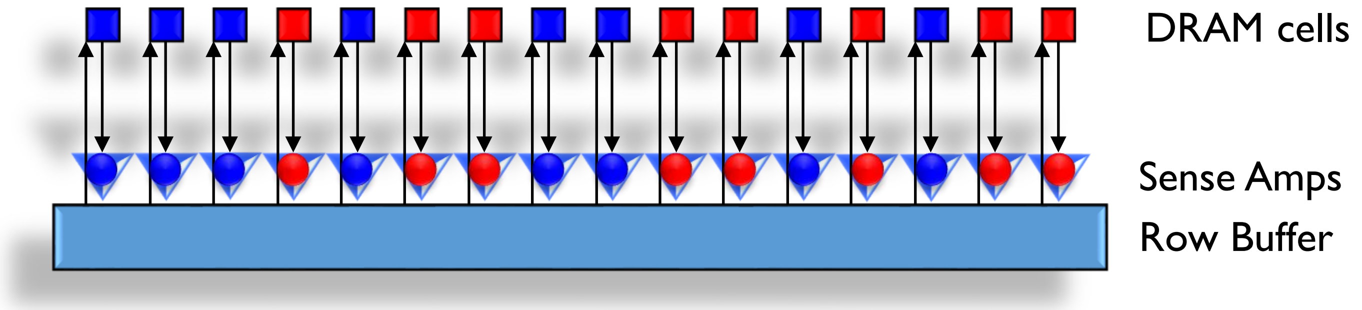
Destructive Read



After read of 0 or 1, cell contents close to ½

DRAM Read

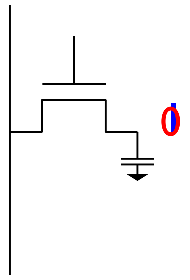
- After a read, the contents of the DRAM cell are gone
  - But still "safe" in the row buffer
- Write bits back before doing another read
- Reading into buffer is slow, but reading buffer is fast
  - Try reading multiple lines from buffer *(row-buffer hit)*



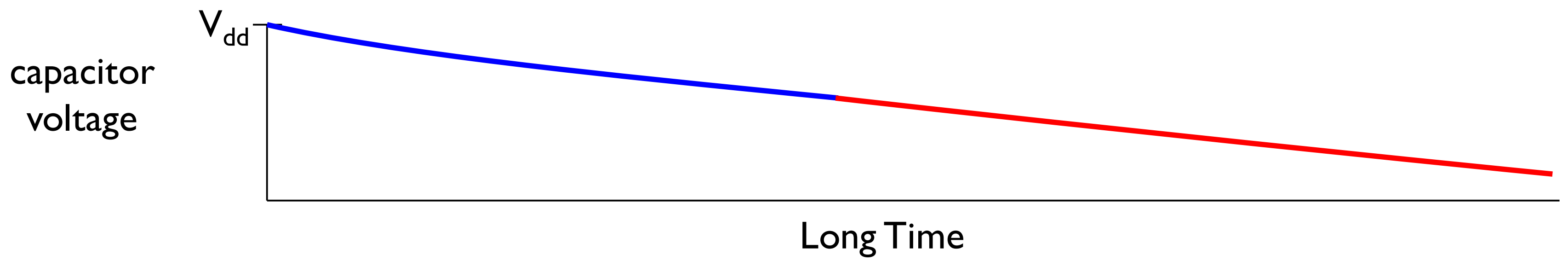
Process is called *opening* or *closing* a row

DRAM Refresh (1/2)

- Gradually, DRAM cell loses contents
  - Even if it's not accessed
  - This is why it's called "dynamic"



- DRAM must be regularly read and re-written
  - What to do if no read/write to row for long time?



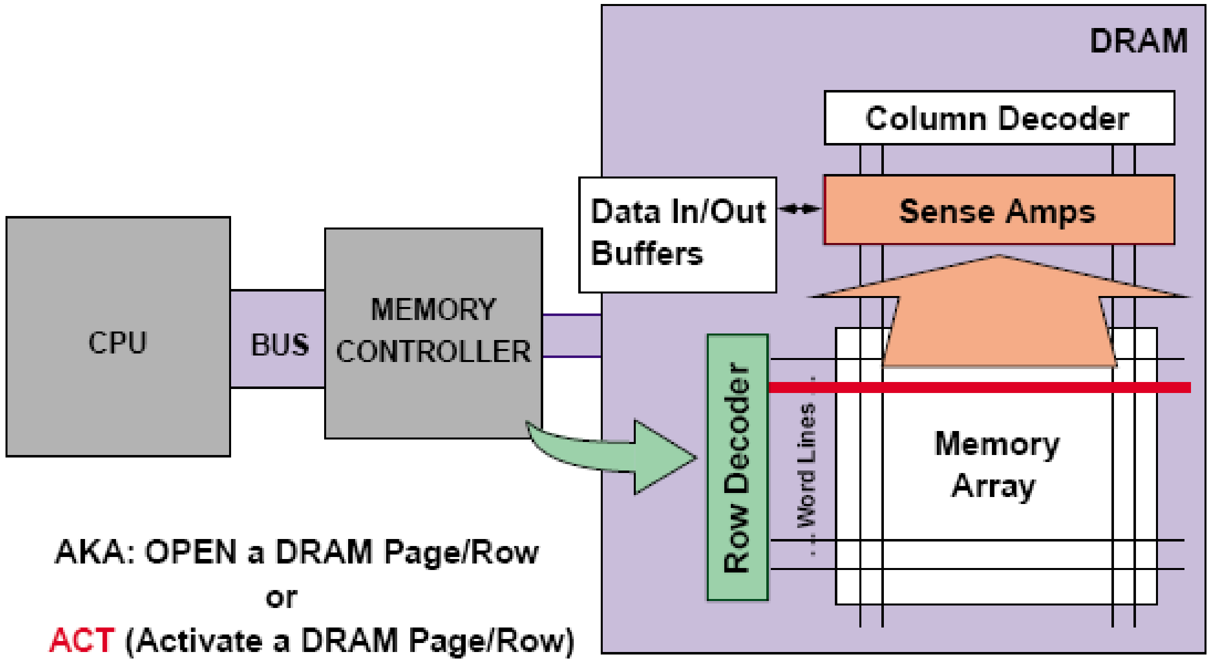
Must periodically *refresh* all contents

DRAM Refresh (2/2)

- Burst Refresh
  - Stop the world, refresh all memory
- Distributed refresh
  - Space out refresh one (or a few) row(s) at a time
  - Avoids blocking memory for a long time
- Self-refresh (low-power mode)
  - Tell DRAM to refresh itself
  - Turn off memory controller
  - Takes some time to exit self-refresh

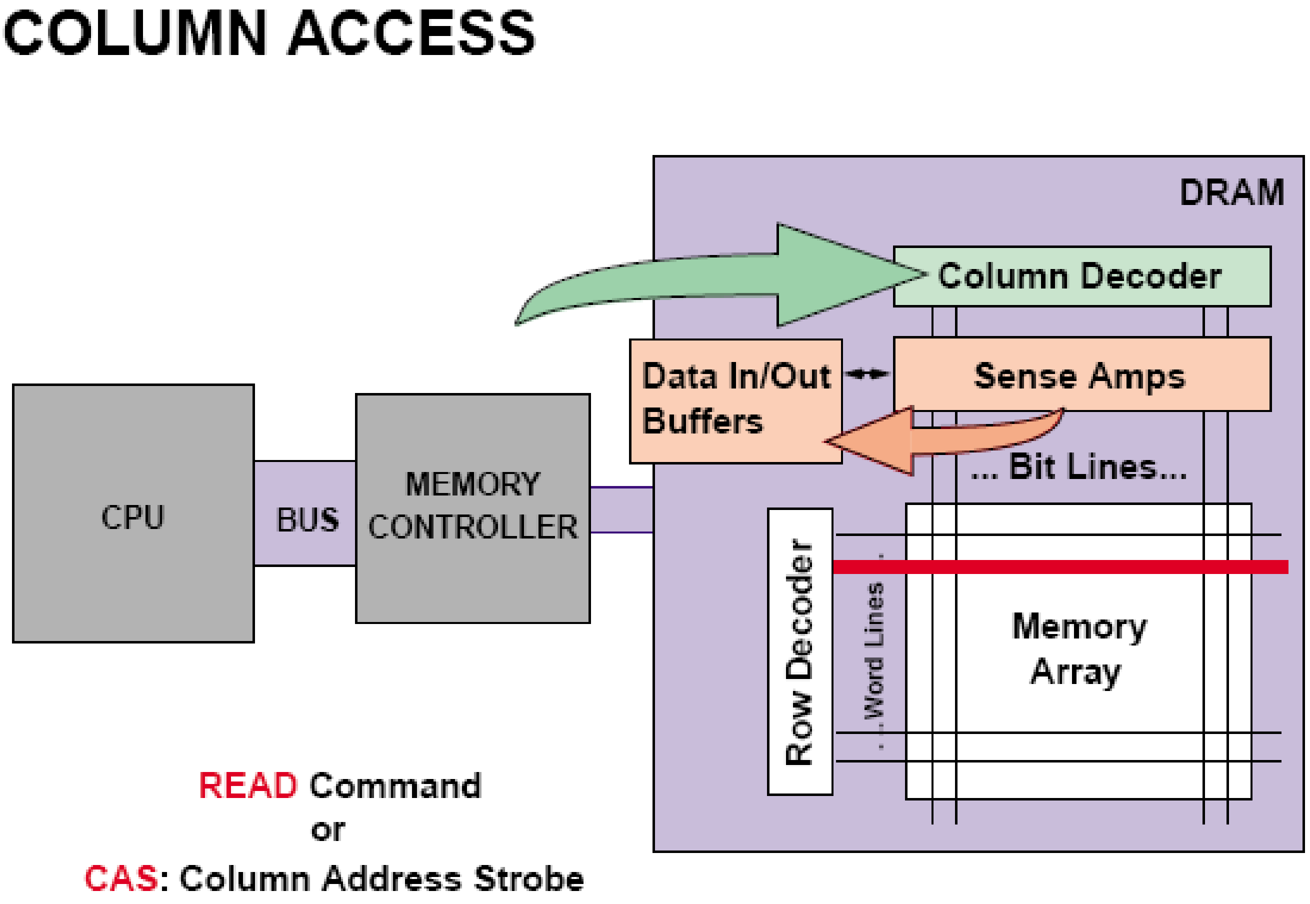
Typical DRAM Access Sequence (1/5)

[PRECHARGE and] ROW ACCESS

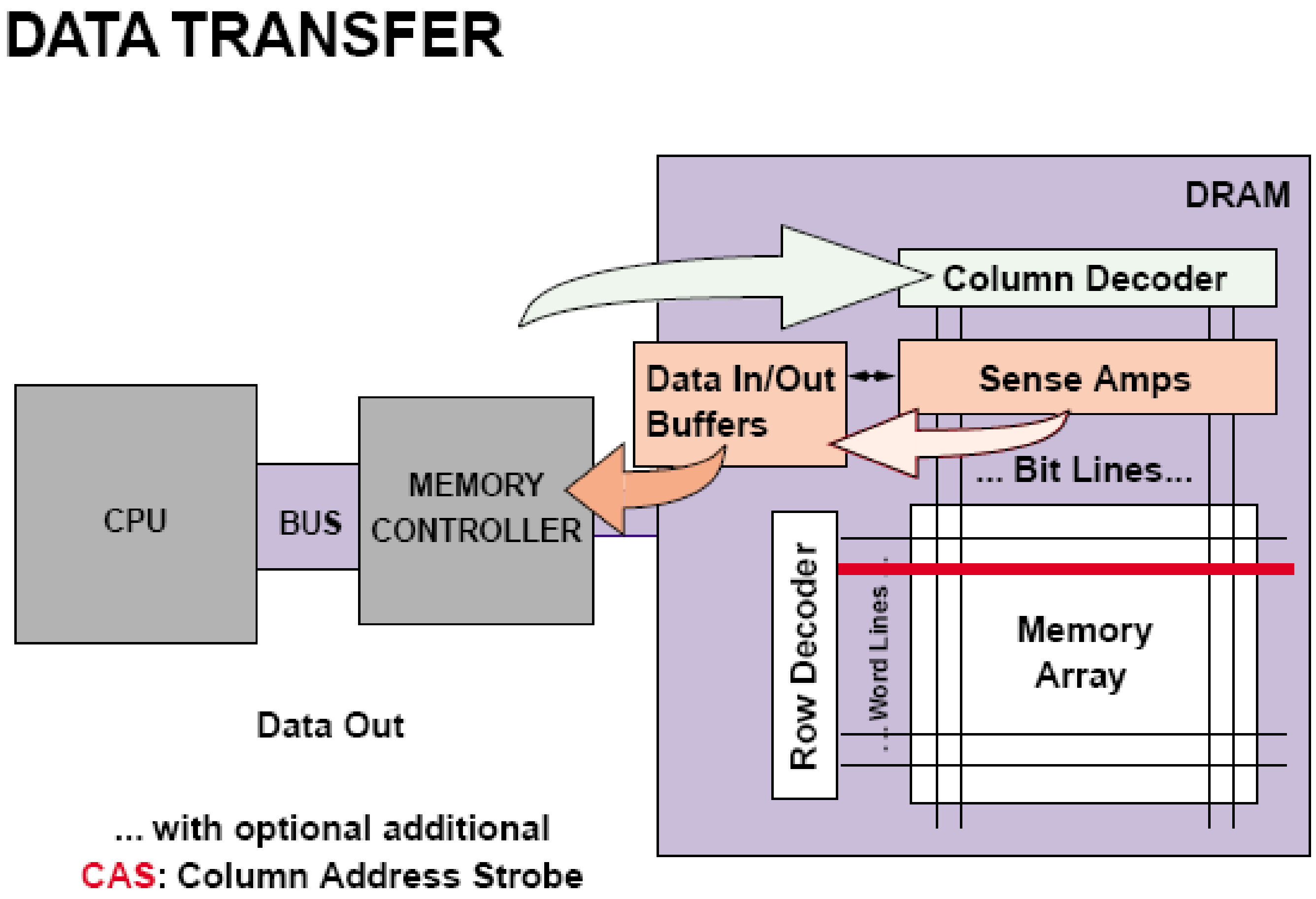


AKA: OPEN a DRAM Page/Row
or
ACT (Activate a DRAM Page/Row)
or
RAS (Row Address Strobe)

Typical DRAM Access Sequence (2/5)

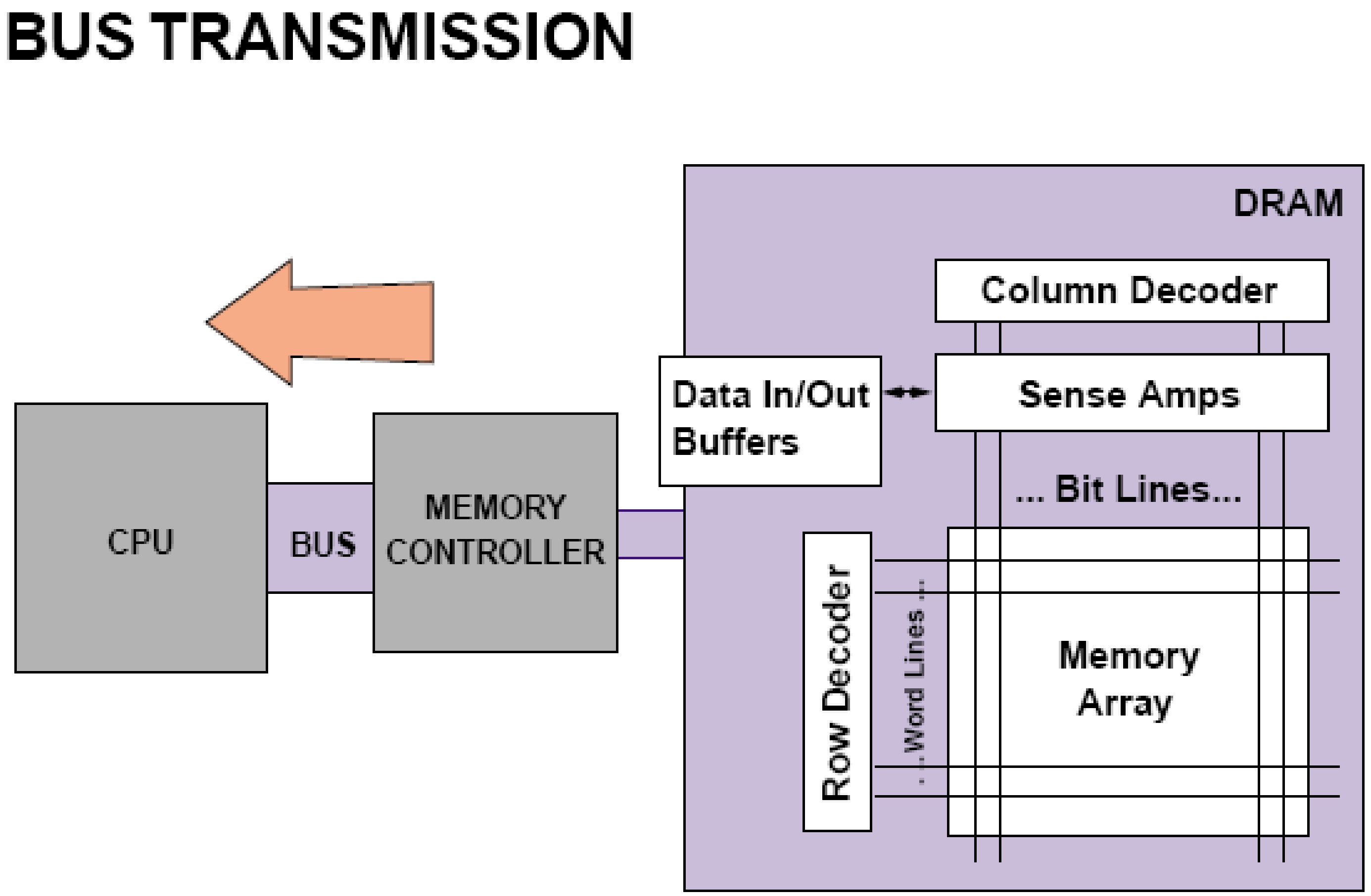


Typical DRAM Access Sequence (3/5)

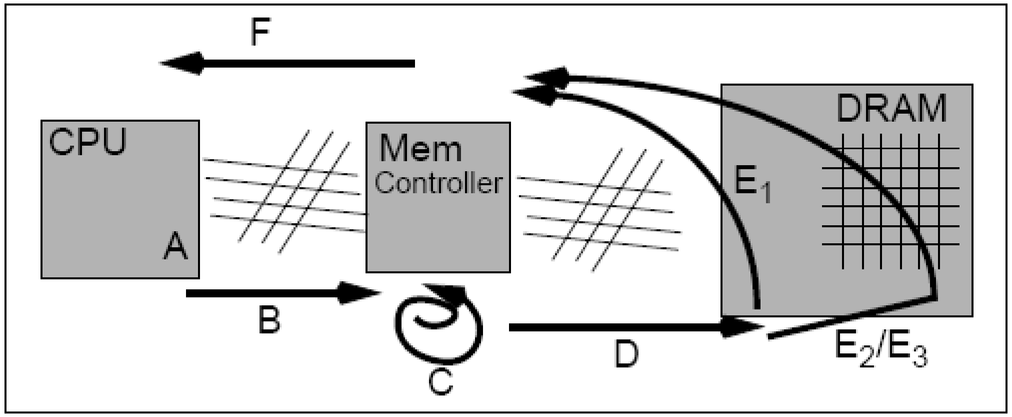


... with optional additional
CAS: Column Address Strobe

Typical DRAM Access Sequence (4/5)



Typical DRAM Access Sequence (5/5)



A: Transaction request may be delayed in Queue
B: Transaction request sent to Memory Controller
C: Transaction converted to Command Sequences
(may be queued)

D: Command/s Sent to DRAM

E_1: Requires only a **CAS** or

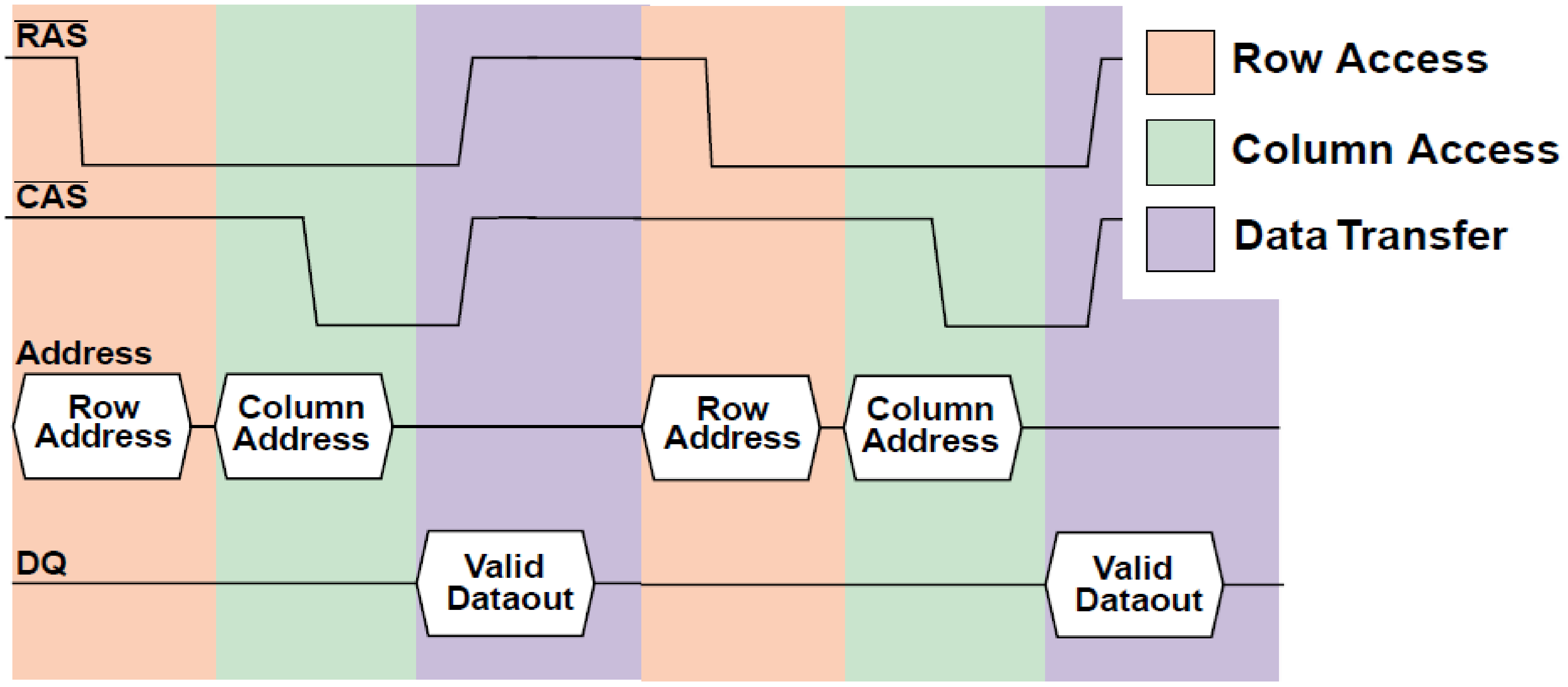
E_2: Requires **RAS + CAS** or

E_3: Requires **PRE + RAS + CAS**

F: Transaction sent back to CPU

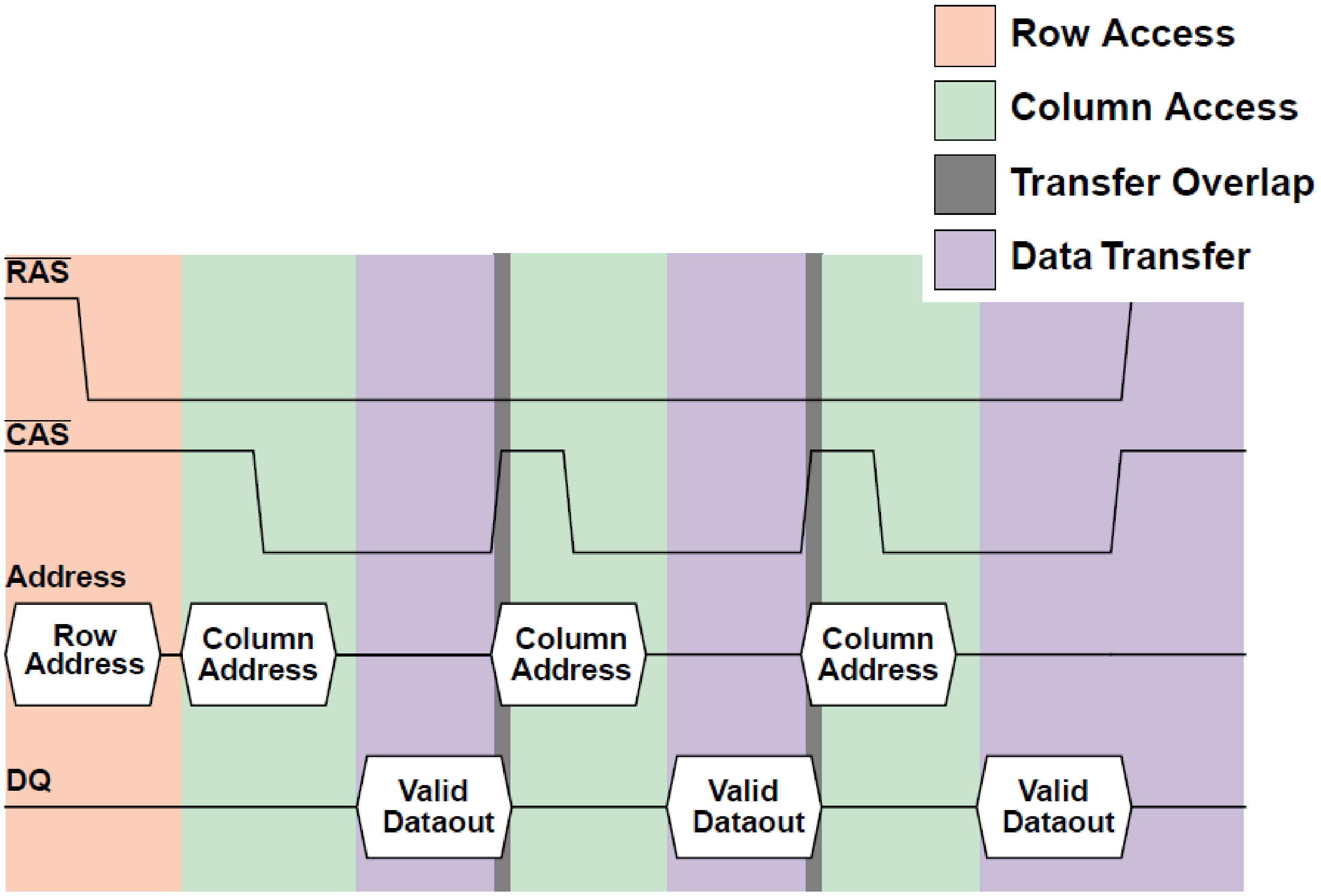
"DRAM Latency" = A + B + C + D + E + F

(Old) DRAM Read Timing



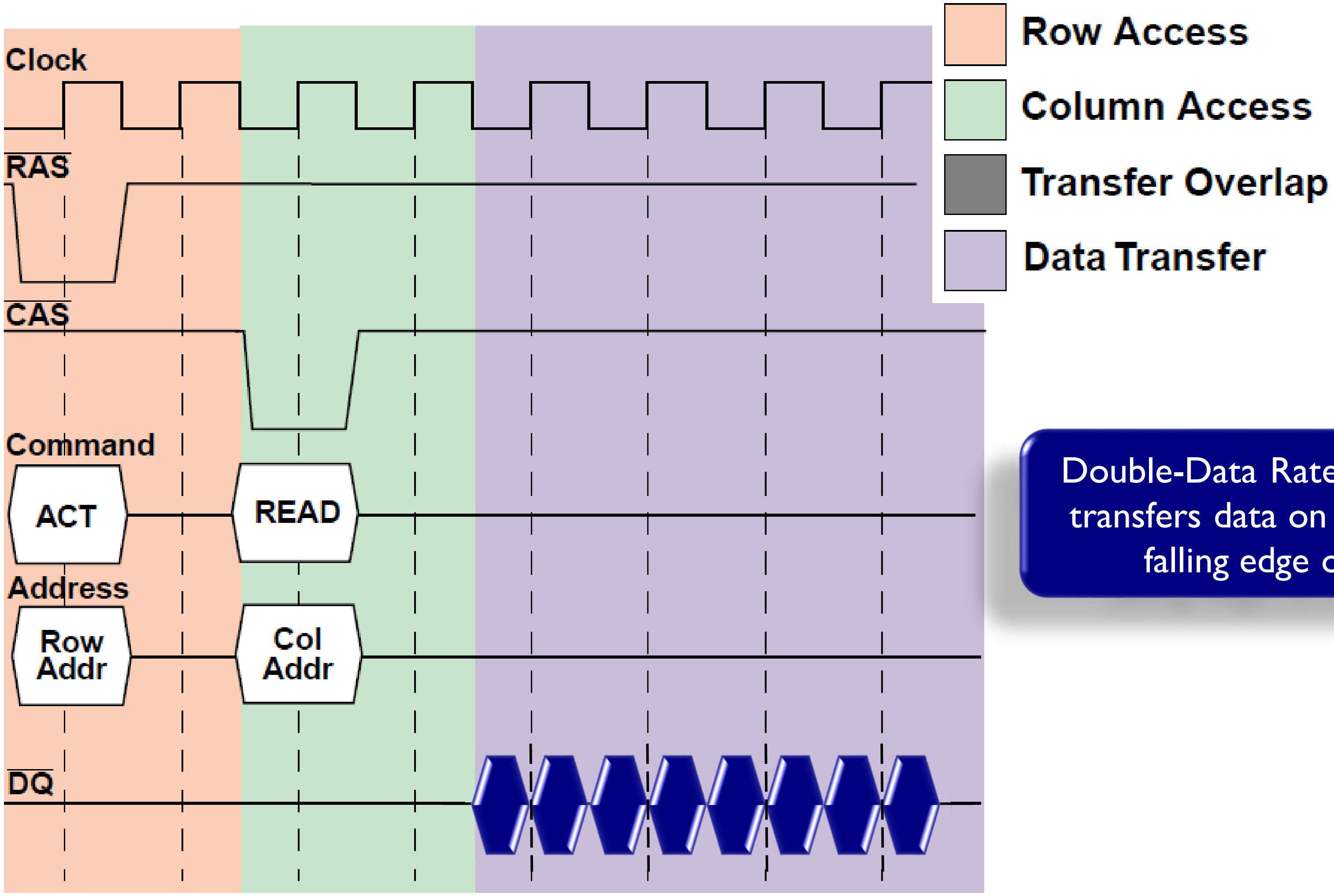
Original DRAM specified Row & Column every time

(Old) DRAM Read Timing w/ Fast-Page Mode



FPM enables multiple reads from page without RAS

(Newer) SDRAM Read Timing



Double-Data Rate (DDR) SDRAM transfers data on **both** rising and falling edge of the clock

SDRAM uses clock, supports bursts

Banking to Improve BW

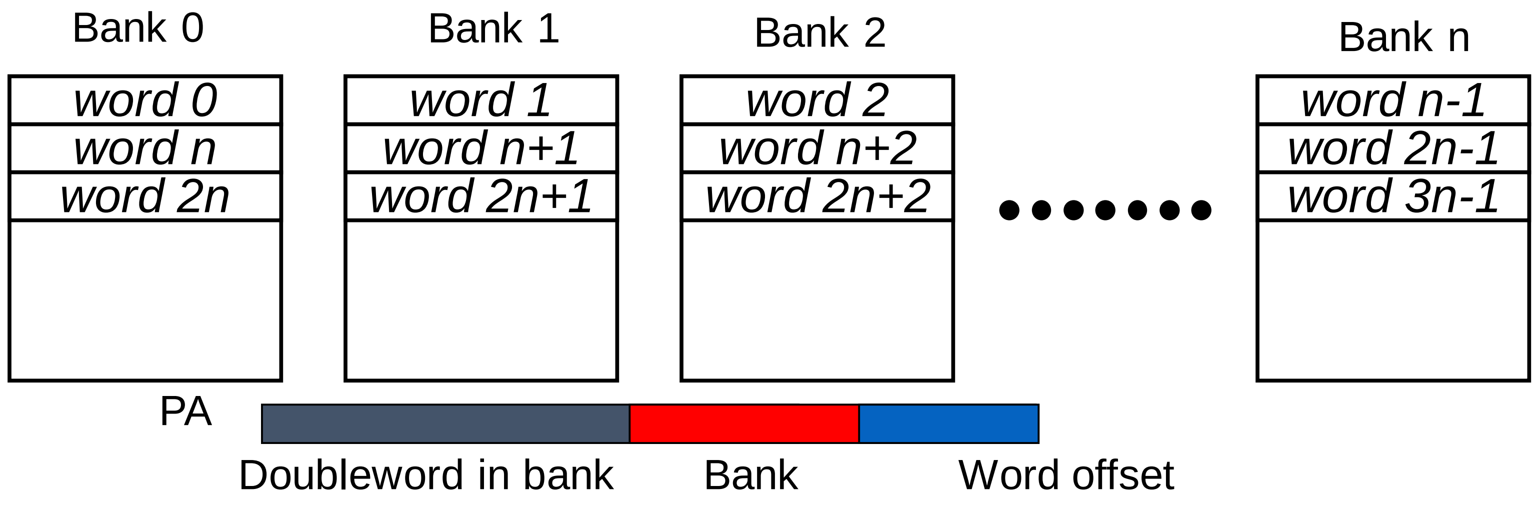
- DRAM access takes multiple cycles
- What is the miss penalty for a 4-word cache block
  - Consider these parameters:
    - 1 cycle to send address
    - 6 cycles to access each word
    - 1 cycle to send word back
  - (1 + 6 + 1) x 4 = 32

How can we speed this up?

- Make memory and bus wider
  - read out all words in parallel
- Miss penalty for 4-word block
  - 1 + 6 + 1 = 8
- Cost
  - wider bus
  - larger expansion size

Simple Interleaved Main Memory

- Divide memory into n banks, "interleave" addresses across them so that word A is
  - in bank (A mod n)
  - at word (A div n)

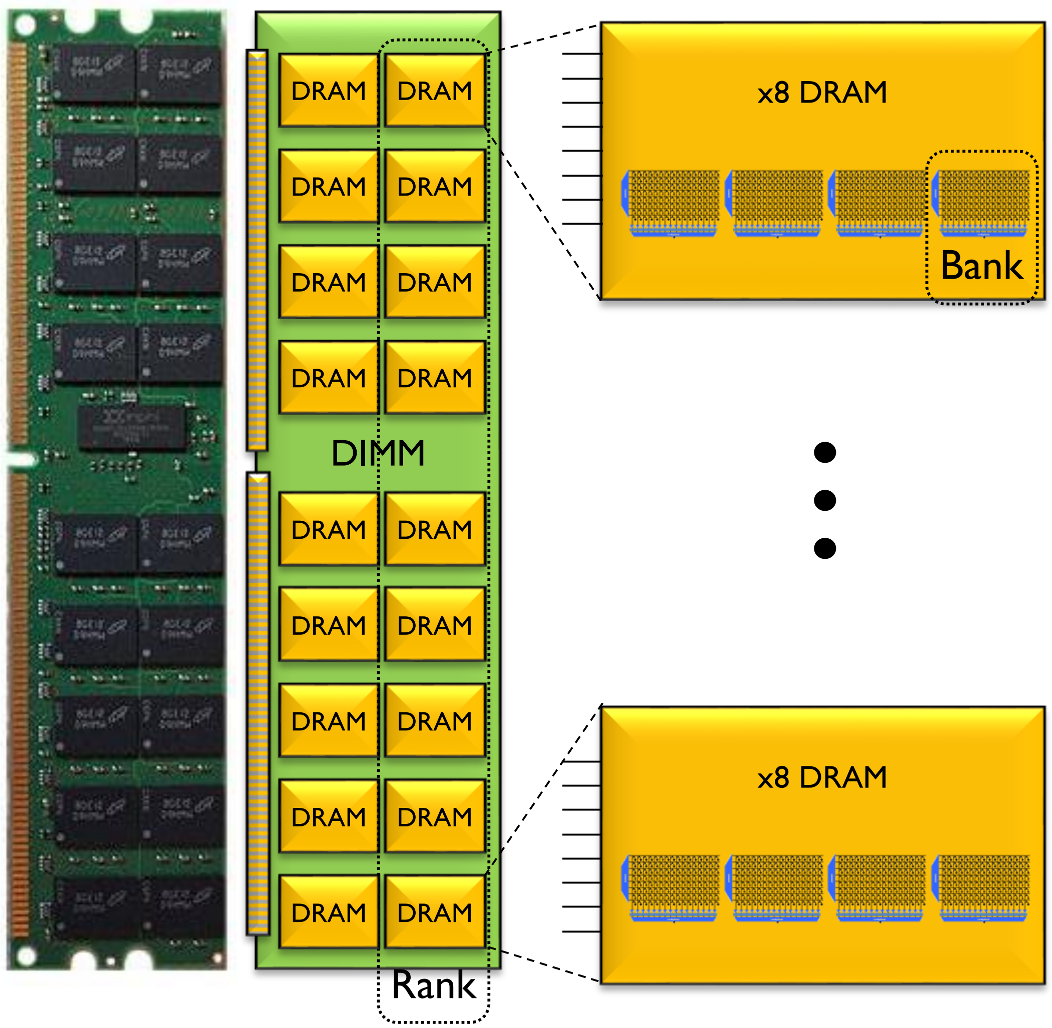


- Can access one bank while another one is busy

→ Interleaving increases memory bandwidth w/o a wider bus

Use parallelism in memory banks to hide memory latency

DRAM Organization



All banks within the rank share all address and control pins

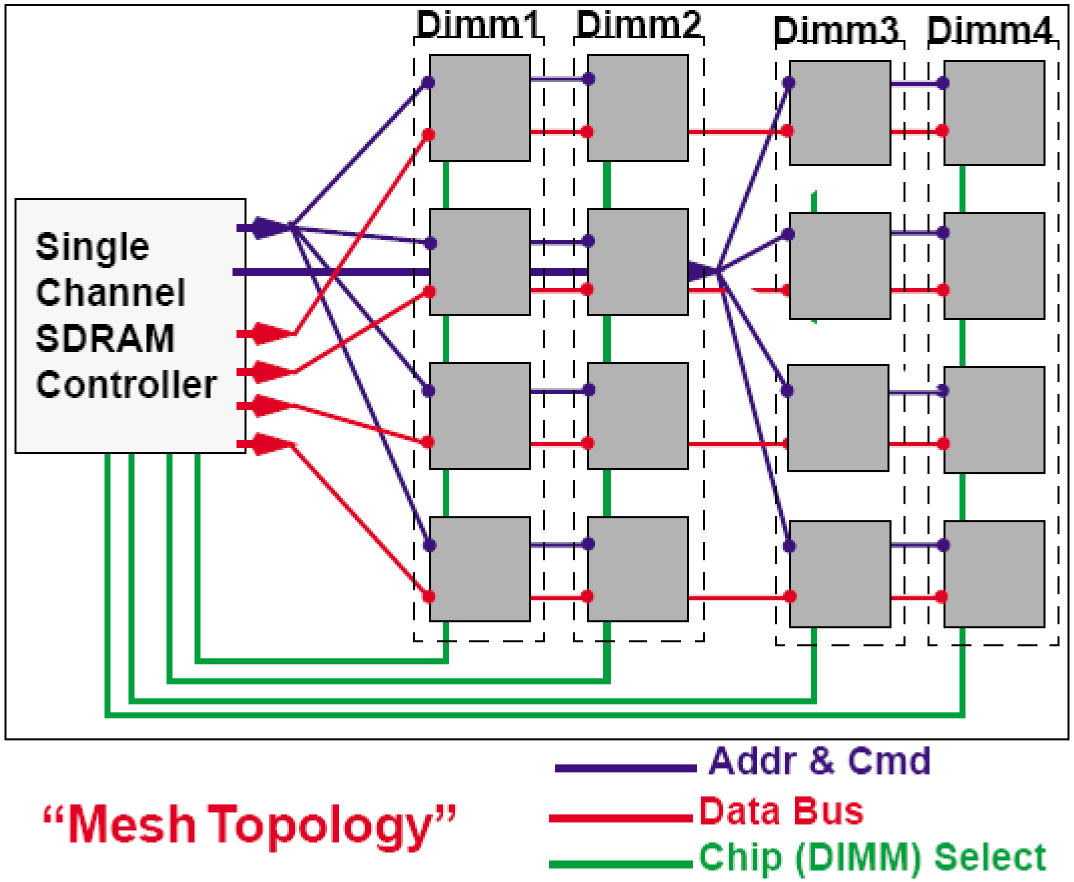
All banks are independent, but can only talk to one bank at a time

x8 means each DRAM outputs 8 bits, need 8 chips for DDRx (64-bit)

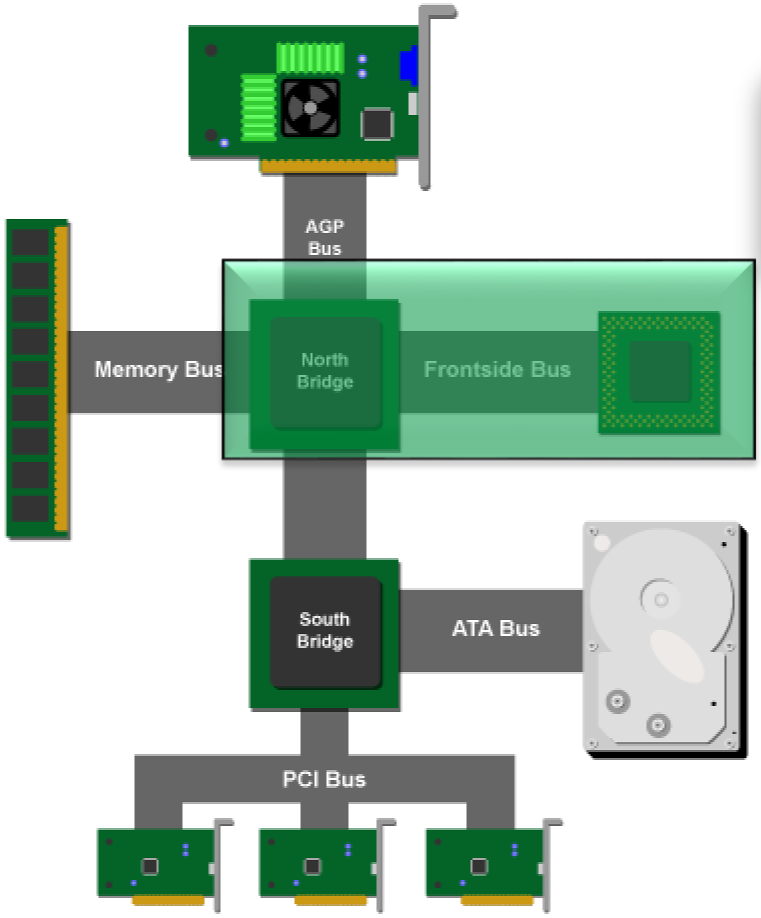
Why 9 chips per rank? 64 bits data, 8 bits ECC

Dual-rank x8 (2Rx8) DIMM

SDRAM Topology



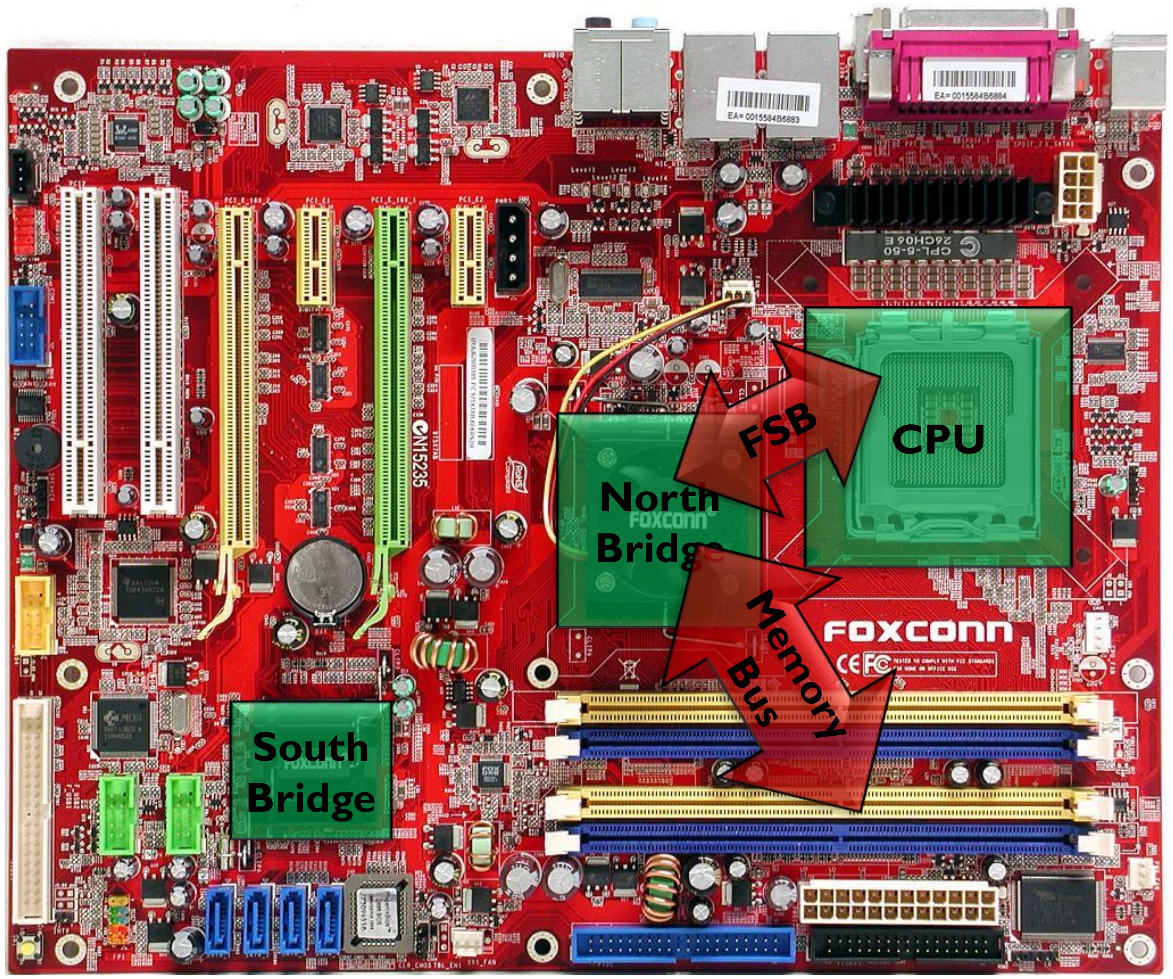
CPU-to-Memory Interconnect (1/3)



North Bridge can be Integrated onto CPU chip to reduce latency

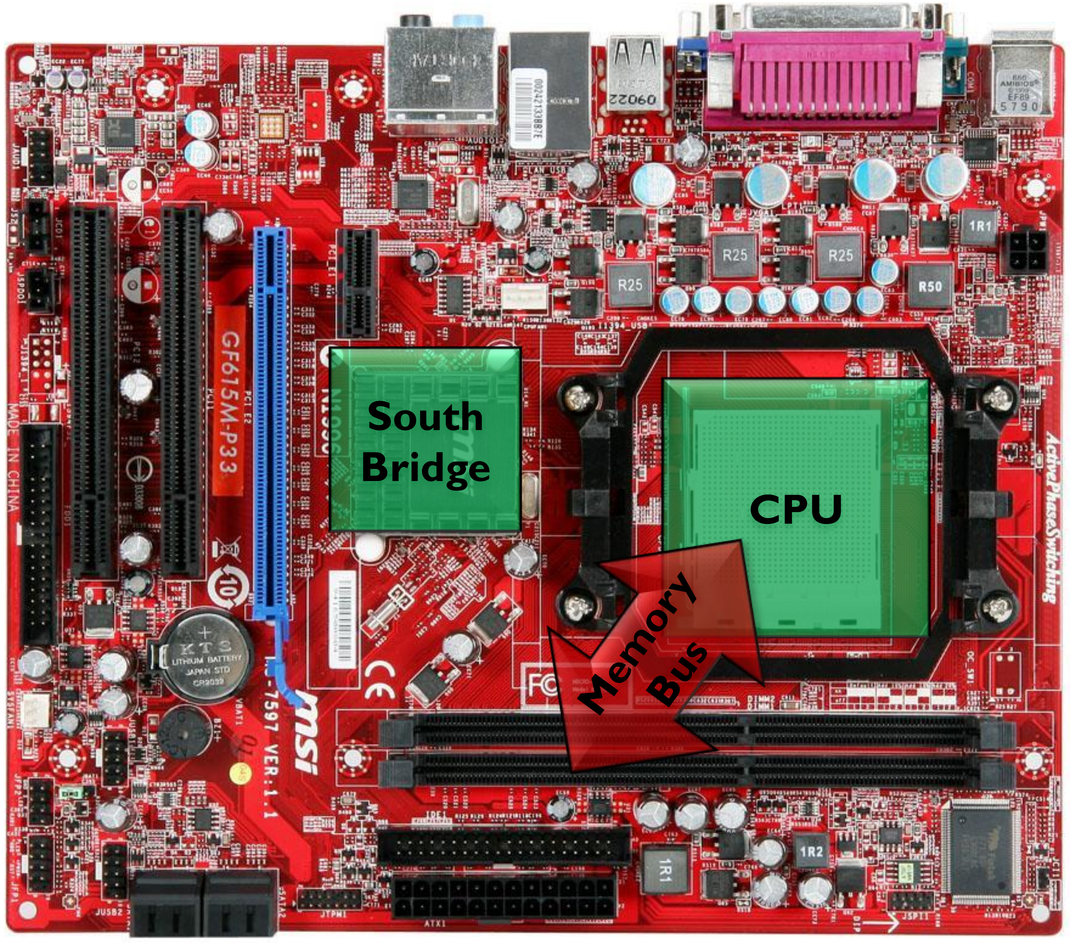
Figure from ArsTechnica

CPU-to-Memory Interconnect (2/3)



Discrete North and South Bridge chips

CPU-to-Memory Interconnect (3/3)



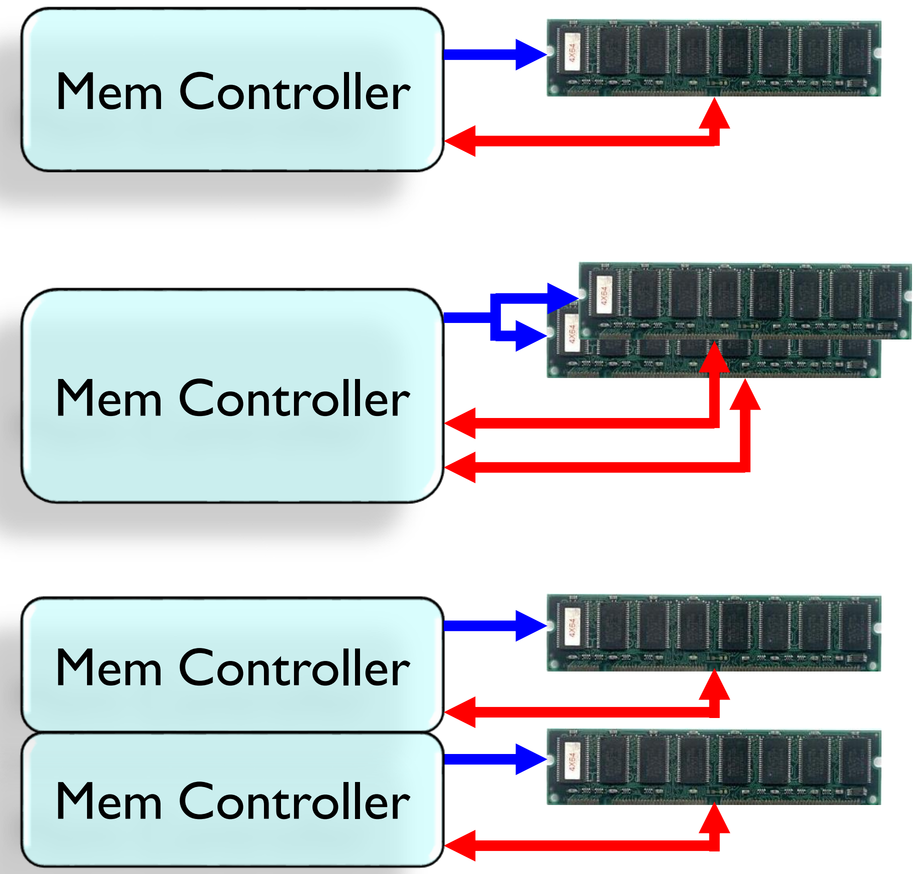
Integrated North Bridge

Memory Channels

One controller
One 64-bit channel

One controller
Two 64-bit channels

Two controllers
Two 64-bit channels



🟦 Commands
🟥 Data

Use multiple channels for more bandwidth

Memory-Level Parallelism (MLP)

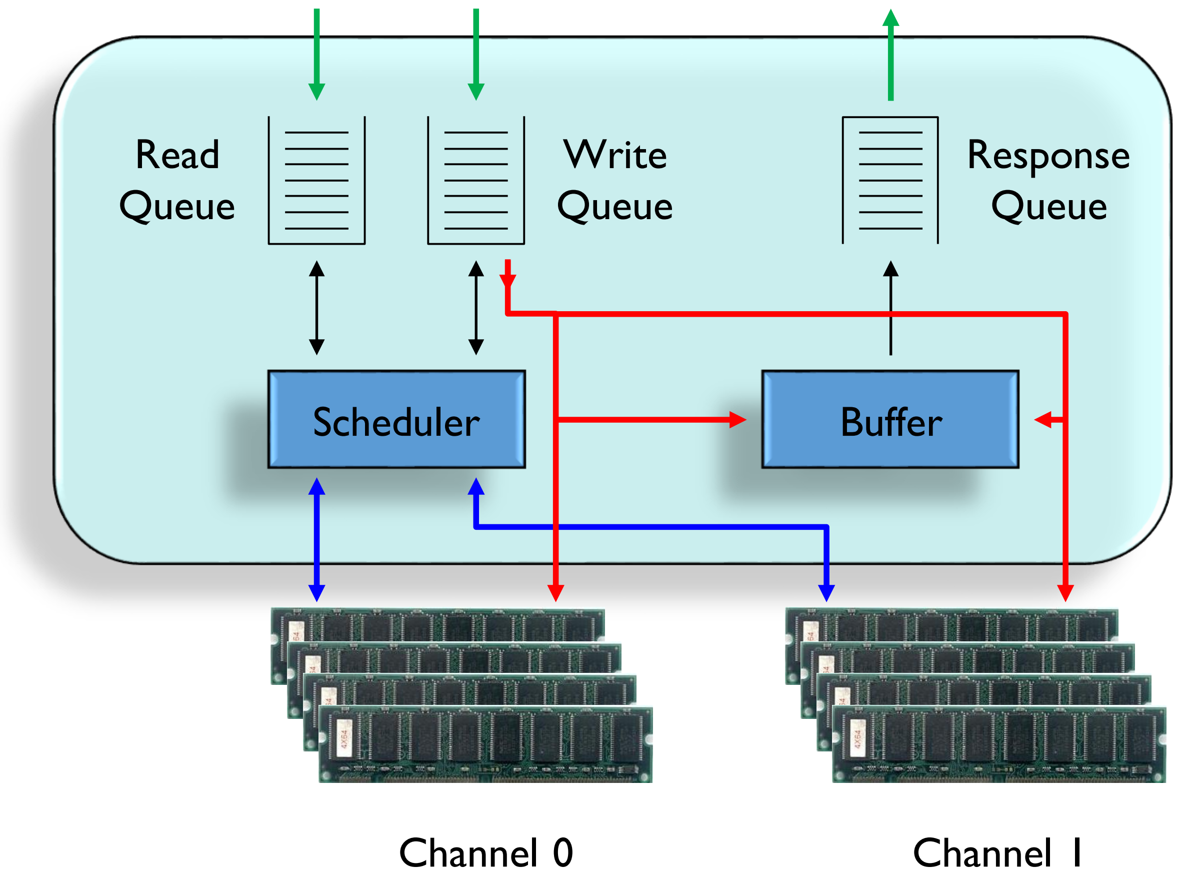
- What if memory latency is 10000 cycles?
  - Runtime dominated by waiting for memory
  - What matters is ***overlapping memory accesses***
- *Memory-Level Parallelism (MLP)*:
  - "Average number of outstanding memory accesses when at least one memory access is outstanding."
- MLP is a metric
  - ***Not*** a fundamental property of workload
  - Dependent on the microarchitecture

AMAT with MLP

- If …
  cache hit is 10 cycles (core to L1 and back)
  memory access is 100 cycles (core to mem and back)
- Then …
  at 50% miss ratio: AMAT = $0.5\times10+0.5\times100 = 55$
- Unless MLP is >1.0, then…
  at 50% mr, 1.5 MLP: AMAT = $(0.5\times10+0.5\times100)/1.5 = 37$
  at 50% mr, 4.0 MLP: AMAT = $(0.5\times10+0.5\times100)/4.0 = 14$

In many cases, MLP dictates performance

Memory Controller (1/2)



🟦 Commands
🟥 Data
🟩 To/From CPU

Memory Controller

Memory Controller (2/2)

- Memory controller connects CPU and DRAM
- Receives requests after cache misses in LLC
  - Possibly originating from multiple cores
- Complicated piece of hardware, handles:
  - DRAM Refresh
  - Row-Buffer Management Policies
  - Address Mapping Schemes
  - Request Scheduling

Request Scheduling (1/3)

- Write buffering
  - Writes can wait until reads are done
- Controller queues DRAM commands
  - Usually into per-bank queues
  - Allows easily reordering ops. meant for same bank
- Common policies:
  - *First-Come-First-Served (FCFS)*
  - *First-Ready—First-Come-First-Served (FR-FCFS)*

Request Scheduling (2/3)

- First-Come-First-Served
  - Oldest request first
- First-Ready—First-Come-First-Served
  - *Prioritize column changes over row changes*
  - *Skip over older conflicting requests*
  - Find row hits (on queued requests)
    - Find oldest
    - If no conflicts with in-progress request → good
    - Otherwise (if conflicts), try next oldest

Request Scheduling (3/3)

- Why is it hard?
- Tons of timing constraints in DRAM
  - tWTR: Min. cycles before read after a write
  - tRC: Min. cycles between consecutive open in bank
  - …
- Simultaneously track resources to prevent conflicts
  - Channels, banks, ranks, data bus, address bus, row buffers
  - Do it for many queued requests at the same time
    … while not forgetting to do refresh

Row-Buffer Management Policies

- *Open-page Policy*
  - After access, keep page in DRAM row buffer
  - Next access to same page → lower latency
  - If access to different page, must close old one first
    - Good if lots of locality

- *Close-page Policy*
  - After access, immediately close page in DRAM row buffer
  - Next access to different page → lower latency
  - If access to different page, old one already closed
    - Good if no locality (random access)

Address Mapping Schemes (1/3)

- Question: How to map a physical addr to channel ID, rank ID, bank ID, row ID, and column ID?
  - Goal: efficiently exploit channel/rank/bank level parallelism
- Multiple ***independent*** channels → max parallelism
  - Map consecutive cache lines to different channels
- Single channel, Multiple ranks/banks → OK parallelism
  - Limited by shared address and/or data pins
  - Map close cache lines to banks within same rank
    - ***Reads*** from same rank are faster than from different ranks
- Accessing different rows from one bank is slowest
  - All requests serialized, regardless of row-buffer mgmt. policies
  - Rows mapped to same bank should avoid spatial locality
- Column mapping depends on row-buffer mgmt. (why?)

Address Mapping Schemes (2/3)

[… … … … bank column …]

0x00000	0x00400	0x00800	0x00C00
0x00100	0x00500	0x00900	0x00D00
0x00200	0x00600	0x00A00	0x00E00
0x00300	0x00700	0x00B00	0x00F00

[… … … … column bank …]

0x00000	0x00100	0x00200	0x00300
0x00400	0x00500	0x00600	0x00700
0x00800	0x00900	0x00A00	0x00B00
0x00C00	0x00D00	0x00E00	0x00F00

Address Mapping Schemes (3/3)

- Example Open-page Mapping Scheme:

High Parallelism:	[row rank bank column channel offset]
Easy Expandability:	[channel rank row bank column offset]

- Example Close-page Mapping Scheme:

High Parallelism:	[row column rank bank channel offset]
Easy Expandability:	[channel rank row column bank offset]

Overcoming Memory Latency

- Caching
  - Reduce average latency by avoiding DRAM altogether
  - Limitations
    - Capacity (programs keep increasing in size)
    - Compulsory misses
- Prefetching
  - Guess what will be accessed next
    - Put in into the cache